## 1 Editor Comment #1

Dear authors,

In my role as Executive editor of GMD, I would like to bring to your attention our Editorial version 1.1: http://www.geosci-model-dev.net/8/3487/2015/gmd-8-3487-2015.html This highlights some requirements of papers published in GMD, which is also available on the GMD website in the Manuscript Types section:

http://www.geoscientific-model-development.net/submission/manuscript\_types.html In particular, please note that for your paper, the following requirement has not been met in the Discussions paper:

• "The main paper must give the model name and version number (or other unique identifier) in the title."

Please add a version number for the ISSM model in the title upon your revised submis- sion to GMD.

Yours,

Astrid Kerkweg

We thank the executive editor for catching this issue, and have accordingly added the ISSM version number to the title.

## 2 Reviewer #1: L. Hascoet

General comments:

The paper describes the construction of the adjoint of the ice sheet simulator ISSM (written in C++) through the use of the overloading-based AD tool ADOL-C, and through the use of the MPI wrapper library AMPI (Adjoinable MPI). The paper presents the first meaningful gradients obtained with this adjoint ISSM, and discusses performance aspects. The main focus of the paper is about the programming strategy that was chosen to ease the algorithmic differentiation process with ADOL-C, the specific efficient tactique to propagate the adjoint through calls to library linear solvers, and the constraints and benefits of using th AMPI library to get an MPI-parallel adjoint. Although not a specialist of the application domain, I have the feeling that the main similar efforts by glaciology researchers are presented and compared adequately. The paper claims that the adjoint ISSM will allow for full exploitation of the enormous quantity of measurement data available. This is something one can believe in. The strategies proposed in sections 2.2 to 2.4 go into details that shed an interesting light on ADOL-C (questions of location handling and contiguity that I was not aware of, thanks) and on the technicalities of the linear solver "trick". Many of these are of general interest for the AD people, overloaders or source-transformers alike. Text is well written, in excellent English. My only minor concern is about its high level of technicity in Computer Science, which might make it hard to read at places, especially for readers specialized rather in geophysics models. But this high level makes it a profitable read for computer science (in particular AD) people. To summarize, I like the paper. I think it is of value for developers of models, and for AD tool developers as well. Getting this value requires effort while reading, though.

We thank Dr. Hascoet for the positive assessment of the manuscript, in particular concerning the applicability of the approach described to the wider Cryosphere community, as well as the AD community. As such, this manuscript is indeed faced with the difficulty of being quite technical. We have therefore strived in this round of edits to clarify, if not simplify, the manuscript where it was possible to do so, without taking away the necessary technicality required for application to the AD community in particular.

Specific comments:

Still, we face this old dilemma between source-transformation AD (probably more efficient) applied to codes in old-fashioned languages (Fortran or C), versus overloading- based AD (slightly heavier in memory) applied to codes in modern languages (nicer to develop). Sentence at line 370 obviously gives the authors opinion on that, but Im wondering if things are really that terrible. After all, in other application domains, people do use source-transformation AD (when available).

The point is well taken, as indeed source-transformation AD in other application domains has gained acceptance. In Cryosphere though, this is far from being the case one could argue. We understand the subjectivity of our point of view on this matter, and have therefore rephrased this part of the conclusion accordingly.

On APMI in this context, the paper might give more hints on what happens while interpreting the trace R(T) i.e. when an AMPI call is registered in T, R(T) calls AMPI primitives at that time too. Otherwise, the reader might wonder how this can work.

Thank you for pointing this issue. We have addressed it by modifying line 240 in the manuscript from "This implies for the adjoint an increment s=s+t, that is, the adjoint is a communication of the adjoint values in the reverse direction." to "This implies for the adjoint an increment s=s+t, that is, the adjoint of a communication is a reversal of the data flow between buffers. Calculating the adjoint of t=s means sending back the value of t and computing s=s+t."

Technical corrections:

- Line 19 "alreay": typo corrected.
- I dont understand the n2 vs n at line 90. Am I missing something? The manuscript from line 80 to line 91 is indeed misleading, and we have addressed this by removing the concept of scaling. The issue is rather how many forward or backward sweeps are necessary to compute the Jacobian components of interest. We have rephrased the paragraph accordingly. This also proved to be confusing to reviewer #2, and should therefore also resolve his concerns.
- Sentence at line 140 look ambiguous to me. Maybe a few commas or small words would help? The sentence is indeed long and convoluted, we broke it down into two simpler sentences.
- Using \* and \*\* in table 2 is unfortunate: guaranteed confusion with the C dereference operator. Agreed, we have replaced \* by +.
- Paragraph at lines 269-275 is hard to follow. Can it be rephrased? We rephrased the paragraph to try and simplify it, in particular, why the MPI\_DataType cannot be used as a reliable value for template parameterization, as it's not well defined at compile time.
- Line 368 "critial". typo corrected.

## 3 Reviewer #2

The paper illustrates how Algorithmic Differentiation (AD) is implemented in the Ice Sheet System Model (ISSM) to compute discrete adjoints and sensitivities. The paper details the operator overloading approach used to achieve AD without disrupting the existing forward capability. Further it explains the parallelization is performed. Overhead factors for computing adjoints are reported for a couple of model runs.

Although the topic of the paper is surely interesting for this journal, I think the exposition is not always clear. In particular, in the introduction there are some typos, inaccuracies and undefined quantities. I would recommend the authors to revise the paper, particularly the introduction, in order to make it more understandable to the reader not used to forward and reverse AD. The paper focuses on several (important) technical details but I think it lacks a mathematical/algorithmic description of the overall framework.

I would recommend the authors to consider a typical sensitivity problem encountered in applications and show how it is solved in ISSM by reverse AD. An exemplar problem could be to find the sensitivity (derivative) of a functional f = f(x, p) with respect to the parameter p, subject to the (nonlinear) constraint c(x, p) = 0 (here f could represent a misfit of some quantity with observational data, and c the ice sheet model). How this problem (or a similar one, maybe time dependent) is solved by ISSM using revers AD? What linear/nonlinear systems need to be solved?

We thank the reviewer for the time spent on the review, and for the legitimate points raised in the manuscript. In particular, the focus on the clarity of the manuscript in the introduction, and the lack of a mathematical framework. The idea of considering a typical sensitivity problem is a propos, as it is exactly what we have shown in Figs 3 and 4. However, we have clearly not conveyed this efficiently. We have therefore reworked the introduction to better explain the AD approach to the non-cognizant scientist, improved or clarified the explanation of the mathematical framework in Eqs. 1 to 3, and more explicitly linked the framework to the examples exposed in Figs. 3 and 4 for North-East Greenland Ice Stream and Upernavik Glacier, which fit exactly the type of non-linear, transient, sensitivity analyses called for by the reviewer as examples.

Some ice-sheet solvers compute sensitivities by solving an additional adjoint equation: BISICLES (Cornford at al. JCP, 2013), Ymir(Isaac et al, SISC, 2015), Albany/FELIX (perego et.al, JGR, 2014 and Tezaur et al. GMD, 2015). Albany/FELIX computes the partial derivatives of the constraint and functional w.r.t. the parameters and the inde- pendent variables using forward AD (Pawlowski et al, Scientific Programming, 2012). A different AD approach is used in Goldberg et al., GMD, 2016. How does the ISSM approach compares to these?

We thank the reviewer for pointing towards this body of litterature. We have included it in the introduction, and compared against our ISSM approach when necessary.

Detailed Review:

- Line 45: How can automatic differentiation be possible when a manual differentiation is not possible? AD is a way to perform differentiation, so if a model is not differentiable then AD would give NaN or Inf.. or similar. The point is well taken. The issue relates more to the feasibility of manual differentiation, which can become exponentially prohibitive. We have rephrased accordingly.
- Line 63: If the expression to be differentiated is implicit than the accuracy of the solution would depend on the implicit solvers and would mostly likely be larger than machine precision. We agree with this point, and have rephrased this sentence to convey the fact that the precision reached by the AD approach is on par with the precision dictacted by the implicit solvers themselves.

- Line 73: Variable y has not been defined. I assume is y = f(x). Thanks for catching this, we have corrected the issue in Eq. 1.
- Eq. 3: I think there is a typo here. Should it be a = a + r; b = b + r? a b There is definitely a typo, thanks for catching it.
- And why r = 0? Please check also expressions at line 81. I think this part should be expanded and the reverse AD method should be presented more clearly. This part has been redesigned, as it was not clear indeed. Here is the new paragraph in the manuscript:

"This is achieved using the partial derivatives of the forward operation  $\phi$  as weights in redistributing the adjoints of the result r to the adjoints of the respective arguments a, b following the rule

$$\bar{a} = \bar{a} + \frac{\partial \phi}{\partial a} \bar{r}; \ \bar{b} = \bar{b} + \frac{\partial \phi}{\partial b} \bar{r}; \ \dots \bar{r} = 0$$
 (1)

Here  $\bar{v}$  denotes the adjoint corresponding to variable v. The final zeroing of the result adjoint  $\bar{r}$  corresponds to the notion that r had been assigned a new value overriding any previous value r might have had. Correspondingly the adjoint of r has been completely distributed and to allow the increment formulation for the adjoint of any prior value of r it has to (re)set to 0."

- Line 82 The number of sweeps depends only on m but the computational cost depends on n as well (as mentioned at line 91). We totally agree. We refer to the same concern raised by Dr. Hascoet, which we already addressed.
- Line 86: In this instance what does f denote? The misfit in the surface elevation? And x? Please describe better in the paper the parallel between the physical problem and the equation. Here, f refers to the function described in Eq. 1. We have corrected the manuscript accordingly. In terms of overall manuscript rewrite, we have now better defined the applications shown in Figs. 3 and 4 (the physical problem) and the description of the problem from Eqs. 1 through 3, as requested by the same reviewer in the general remarks also.
- Line 201: Hows b is defined? Please expand a bit this part to help the reader better understand how an implicit equation as As = r is handled by reverse AD. Also, in the introduction r was used to denote an elemental operation (or its output), whether here it is used to denote the residual of the system, which is an input of the el- emental" operation s = A1r, I think this can be confusing. Lets say only a few elements of A are active (say, the part related to the boundary, if A comes from discretizing a PDE), then do you need to store the entire matrix A as well? What happens when the operator A is non linear (it depends on s)? There is a typo, and b is now replaced by r. Indeed the choice of r vs b is unfortunate, and

we have reverted to b by default here. In terms of operation, the As = r operation represents an elementary operation, and does not represent a non-linear system of equations per se. This also mitigates the question from the reviewer regarding boundary conditions: the matrix A here is the matrix corresponding to the free degrees of freedom, after reduction from the global set of degrees of freedom by constraining the degrees of freedom lying on the Dirichlet boundary conditions of the system. But it could equally well correspond to any equation system As=r where A is not a stiffness matrix (in Finite Element Modeling). We have rewored the manuscript accordingly here to make this point clear.

- Line 324 Again, I dont think that the computational cost do not depend on n, even theoretically.Agreed, we have rephrased to convey the fact that the computation relies on one reverse sweep, hence the need to compare the computation time between one forward sweep to evaluate f, and one reverse sweep to evaluate the gradient of f.
- Figure 5-6 What is the specific application solved to produce results showed in Figures4and 5? Are these non linear problems? How many processors are you using when solving the probem with MUMPs? Please add also the overloading overhead in Figure 6, as done in Figure 5. The manuscript has been modified with respect to Figures 3 and 4 to demonstrate the applicability of the AD approach to computing the gradient of a function with respect to model inputs, as requested by the reviewer in his general remarks also. We have also added a description of the benchmark test relied upon in the text, corresponding to test 101 of the ISSM regression suite, and details about the number of cpus used with the MUMPS problem. In terms of non-linearity, test 101 is non-linear with respect to the way viscosity is treated, we have also described this further in the manuscript.

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# An Approach to Computing Discrete Adjoints for MPI-Parallelized Models Applied to the Ice Sheet System Model 4.11

Eric Larour<sup>1</sup>, Jean Utke<sup>2</sup>, Anton Bovin<sup>3</sup>, Mathieu Morlighem<sup>4</sup>, and Gilberto Perez<sup>5</sup>

<sup>1</sup>Jet Propulsion Laboratory - California Institute of technology, 4800 Oak Grove Drive MS 300-323, Pasadena, CA 91109-8099, USA
<sup>2</sup>Allstate Insurance Company, 2775 Sanders Rd., Northbrook, IL 60062, USA.
<sup>3</sup>The University of Chicago, Chicago, Illinois, USA.
<sup>4</sup>University of California, Irvine, Department of Earth System Science, Croul Hall, Irvine, CA 92697-3100, USA
<sup>5</sup>University of California, Irvine, School of Information and Computer Sciences, Irvine, CA 92697-3100, USA

Correspondence to: Eric Larour (eric.larour@jpl.nasa.gov)

### Abstract.

Within the framework of sea-level rise projections, there is a strong need for hindcast validation of the evolution of polar ice sheets in a way that tightly matches observational records (from radar, gravity, and altimetry observations mainly). However, the computational requirements for making

- 5 hindcast reconstructions possible are severe and rely mainly on the evaluation of the adjoint state of transient ice-flow models. Here, we look at the computation of adjoints in the context of the NASA/JPL/UCI Ice Sheet System Model, written in C++ and designed for parallel execution with MPI. We present the adaptations required in the way the software is designed and written but also generic adaptations in the tools facilitating the adjoint computations. We concentrate on the use of
- 10 operator overloading coupled with the AdjoinableMPI library to achieve the adjoint computation of ISSM. We present a comprehensive approach to 1) carry out type changing through ISSM, hence facilitating operator overloading, 2) bind to external solvers such as MUMPS and GSL-LU and 3) handle MPI-based parallelism to scale the capability. We demonstrate the success of the approach by computing sensitivities of hindcast metrics such as the misfit to observed records of surface altime-
- 15 try on the North-East Greenland Ice Stream, or the misfit to observed records of surface velocities on Upernavik Glacier, Central West Greenland. We also provide metrics for the scalability of the approach, and the expected performance. This approach has the potential of enabling a new generation of hindcast-validated projections that make full use of the wealth of datasets currently being collected, or alreay already collected in Greenland and Antarctica.

#### 20 1 Introduction

Constant monitoring of polar ice sheets through remote sensing, in particular since the advent of altimeter, radar, and gravity sensors such as ICESat-1, CryoSat, RADARSAT-1, ERS-1 and ERS-2, Envisat, and GRACE, has created a large amount of data that has yet to find its way through Ice Sheet Models (ISMs) and hindcast reconstructions of polar ice sheet evolution. In particular, as

- 25 evidenced by the wide discrepancy between ISMs involved in the SeaRISE and Ice2Sea projects (Nowicki et al., 2013; Bindschadler et al., 2013) significant improvements in modeled projections of the mass balance of polar ice sheets and their contribution to future sea-level rise has not resulted from the increase in availability of data, but rather from improvements in the type of physics captured in forward models. One reason for this is the lack of data assimilation capabilities embedded in
- 30 the current generation of state-of-the-art ISMs. In the past 10 years, great strides have been made in improving model initialization by using steady-state model inversions of basal friction (MacAyeal, 1993; Morlighem et al., 2010; Larour et al., 2012; Price et al., 2011; Arthern and Gudmundsson, 2010), ice rheology (Rommelaere and MacAyeal, 1997; Larour et al., 2005; Khazendar et al., 2007) and bedrock elevation (Morlighem et al., 2014) among others. However, these approaches aim at
- 35 improving our knowledge of poorly constrained input parameters and boundary conditions as long as the ice-flow regime is captured in a steady-state configuration. These inversions rely on analytically derived adjoint states of forward stress-balance or mass-transport models, but do not extend to transient regimes of ice flow.

Applications to transient models and long temporal time series such as the ICESat/CryoSat con-

- 40 tinuous altimetry record from 2003 to present-day, have been much more rare, and to our knowledge, are limited to a few studies such as Heimbach and Bugnion (2009), Goldberg and Sergienko (2011), Goldberg and Heimbach (2013), Larour et al. (2014) and Goldberg et al. (2015) among others. The main issue here precluding widespread application of transient data assimilation lies in the difficulty of deriving temporal adjoints of transient models. Actually, in In many cases, a the sheer amount
- 45 of physical processes being represented makes the manual derivation of the adjoint state of a forward model is not possible, especially where ice-flow physics are not differentiable. This is the case for example in thermal transient models, where the melting-point is a physical constraint (of the threshold type) that is imposed on temperature. This numerical issue extremely cumbersome and difficult to mantain. This state of affairs can be mitigated by adopting different approaches, such
- 50 as: 1) ensemble runs, as in Applegate et al. (2012), where model runs compatible with observations are selected; 2) methods similar to the flux-correction methods implemented in Aschwanden et al. (2013); Price et al. (2011) where boundary conditions are corrected in order to match time series of observations (tuning approach); 3) quasi-static approaches, where snapshot inversions are carried out in time, as in Habermann et al. (2012, 2013) and 4) sampling methods, which have the main
- 55 drawback of being computationally very expensive (each sample at the cost of one forward run). Though this is not an exhaustive list of all available methods, the main advantage of adjoint driven

inversions is that it relies on the exact sensitivity of a forward model to its inputs, hence ensuring a physically meaningful inversion.

Understanding sensitivities of a forward ice-flow model, which is needed to physically con-60 strain a temporal inversion, requires computation of derivatives of model outputs to model inputs. If such derivatives are approximated by finite-difference schemes, they are subject to the tradeoff between approximation and truncation errors for the perturbation, which is aggravated for higherorder derivatives. If the derivatives are computed using algorithmic differentiation (AD) (Griewank and Walther, 2008), also known as automatic differentiation, then one can attain derivatives with

65 machine precision provided at a floating point precision equivalent to that of the underlying program implementation of the numerical model, provided it is amenable to the application of an AD tool. In particular, this approach does not depend on the type of physics relied upon, and it is transparent to the model equations, provided each step of the overall software is differentiable. Indeed, the AD method assumes a numerical model *M* that is a function

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$$\begin{array}{cccc} f: \mathbb{R}^n & \to & \mathbb{R}^m & f: \mathbb{R}^n & \to & \mathbb{R}^m \\ x & \mapsto & f(x) & x & \mapsto & y = f(x) \end{array}$$
(1)

implemented as a computer program. The execution of the program implies the execution of a sequence of arithmetic operators such as +, -, \* and intrinsics  $\sin, e^x$ , and so forth to which the chain rule can be applied. For each such elemental operation  $r = \phi(a, b, ...)$  with result r and arguments  $a, b, ...^1$  we can write the total derivative as:

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$$\dot{r} = \frac{\partial \phi}{\partial a} \dot{a} + \frac{\partial \phi}{\partial b} \dot{b} + \dots$$
 (2)

For example, if  $\phi$  is the multiplication operator \*, i.e. r = ab then one will get the product rule  $\dot{r} = b\dot{a} + a\dot{b}$ . Applying the above rule to each elemental operation in the sequence gives a method to compute:

$$\dot{\mathbf{y}} = \mathbf{J}\dot{\mathbf{x}}$$
 with the Jacobian  $\mathbf{J} = \begin{bmatrix} \frac{\partial f_i}{\partial x_j} \end{bmatrix}, i = 1 \dots m, \ j = 1 \dots n$ 

without explicitly forming the Jacobian. This method applies the chain rule in the computation order of the values in the program and is known as *forward*-mode AD. The opposite order of applying the chain rule to the elemental operations, known as *reverse* or *adjoint*-mode AD, yields projections  $\bar{x} = J^T \bar{y}$ .

<sup>&</sup>lt;sup>1</sup> In practice most  $\phi$  are uni- or bivariate.

80 This is achieved by applying to the elemental operations using the partial derivatives of the forward operation  $\phi$ , in reverse order of their original execution, the rule: as weights in redistributing the adjoints of the result *r* to the adjoints of the respective arguments *a*, *b* following the rule

$$\bar{a} = \bar{a} + \frac{\partial \phi}{\partial \bar{a}} \frac{\partial \phi}{\partial a} \bar{r}; \ \bar{b} = \bar{b} + \frac{\partial \phi}{\partial b} \bar{r}; \ \dots \bar{r} = 0 \quad .$$
(3)

where the bar operator defines the corresponding adjoint variable. Here v denotes the adjoint
85 corresponding to variable v. The final zeroing of the result adjoint r corresponds to the notion that r had been assigned a new value overriding any previous value r might have had. Correspondingly the adjoint of r has been completely distributed and to allow the increment formulation for the adjoint of any prior value of r it has to (re)set to 0.

For the multiplication example one therefore gets a = a + br; b = b + ar; r = 0. In particular,
90 for applications in which m ≪ n, the reverse mode is advantageous because its computational cost does not depend on depends on m reverse sweeps instead of n forward sweeps. Typical problems in Cryosphere science involve computations of diagnostics which are scalar-valued cost functions (m = 1), such as for example the spatio-temporally averaged misfit between modeled surface elevation and observed surface topography (Larour et al., 2014). For these cases, one can compute the gradient

- 95  $\nabla f = \mathbf{J}^T \bar{y}$  with  $\bar{y} = 1$  as a single projection —(where *f* is the function defined in Eq. 1). Thus, for high-resolution models implying very large *n*, the reverse mode is an enabling and potentially very efficient technique. This significant capability in AD is what makes its application to data assimilation so efficient. Instead of evaluating a Jacobian using one forward sweep for each one of the outputs of the forward model with respect to each input model inputs, which would be significantly
- 100 consuming as it scales in  $n^2$ , the reverse mode evaluates the gradient of one specific output of interest with respect to the model inputs in only one sweep that scales in n. reverse sweep.

Using this approach, it is possible, without manual derivation, to evaluate the derivative  $\frac{\partial f}{\partial x}$  from Eq. 1. In Cryosphere models, f could typically be a transient ice flow model, with inputs x taken for example as the basal drag coefficient  $\alpha$  at the ice/bed interface, and the output y as the cost function

- 105 J measuring the spatio-temporally averaged difference between for example surface elevation and an altimetry record. In this case, m = 1, and n is the number of degrees of freedom of the basal drag input. Obtaining the sensitivity function  $\frac{\partial f}{\partial x}$  is critical both in terms of enabling data assimilation using inversion methods, but also in terms of understanding the sensitivity of a model to its model inputs, as well as the impact of processes represented in the model itself. Other approaches than
- AD can be relied upon to compute such a sensitivity, such as solving for the adjoint equation directly (Morlighem et al., 2010; Larour et al., 2012; Cornford et al., 2013), but usually, a mix of forward AD or reverse AD is implemented for formulations where the manual derivation is difficult (Pawlowski et al., 2012; Goldberg et al., 2016). In Pawlowski et al. (2012) for example, AD relies on template based generic programming to compute sensitivities, while in Goldberg et al. (2016).

115 AD is based on source to source transformation. Our approach will instead be based on so called operator overloading, which is uniquely suited to a framework written in C++.

Applying AD to large-scale frameworks such as ISSM (Larour et al., 2012), MITgm (Heimbach, 2008), SICOPOLIS (Heimbach and Bugnion, 2009) or DassFlow (Monnier, 2010) is a difficult proposition, but one which enables significant improvements in the way models can be initialized

- 120 (Heimbach and Bugnion, 2009), hindcast validated (Larour et al., 2012), and calibrated (Goldberg et al., 2015) towards better projections. Traditional approaches relying on Source-to-Source transformation have been developed, but for frameworks such as ISSM, which are C++ oriented, and highly parallelized, this type of approach breaks down. Our goal here is to demonstrate how the so-called operator-overloading AD approach can be implemented and validated for a framework such
- 125 as ISSM, and what developments were necessary to make this capability operational. Our approach is discussed in section 2 of this manuscript, with section 3 describing the method validation as well as applications with ISSM. We discuss and conclude in the last section the applicability of such an approach to other frameworks, and on the opportunities these new developments afford for Cryosphere Science and data assimilation of remote sensing data in particular.

#### 130 2 Methodology and Validation

#### 2.1 Source to Source vs Overloaded Operators

Distinctions among the AD implementation approaches relate to the way the derivative data in  $\dot{v}$  or  $\bar{v}$  respectively is associated with the original program variable v (data augmentation) and how the logic for coefficient propagation is added to the original program (logic augmentation). The basic options

- 135 for the former are *association by address* and *association by name*. Association by address packs the original and the derivative data into a new type called the *active* type, and all differentiated program variables have their type changed to that new active type. For the logic augmentation one uses *source code transformation* or *operator overloading*. Because of the complexity of the C++ syntax and semantics, there currently still is no comprehensive AD tool for source code transformation of C++
- 140 models and thus operator overloading remains the method of choice, implying association by address for the data augmentation. In practice, the overloaded operators and intrinsics for the forward mode typically execute Eq. (2) directly for the forward mode and for the reverse mode record the sequence of operations into a trace  $T(f(\mathbf{x}))$  which then is read backward by an interpreter R(T) that executes the statements Eq. (3) for the reverse mode. The invocation of R(T) is part of the user-written logic
- 145 that then uses the derivatives. The model thus enhanced by an adjoint capability computing both f and  $\nabla f$  is denoted by  $\overline{M}$ .

#### 2.2 Type Change to Enable Overloaded Operators

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Changing to the aforementioned active type is a significant effort to be undertaken in the model code. Among the choices to effect this type change one should select one that is transparent to the model development process, is maintainable, and minimizes the manual effort.

- The tool of choice for this paper is ADOL-C (Griewank et al., 1996; ADOL-C) and our target model M is the Ice Sheet System Model (ISSM, Larour et al. (2012)). The ADOL-C manual is quiet on the practical approach to effect the type change. For models with a large C++ code base and many contributing developers such as ISSM, not all of them aware of  $\overline{M}$  but merely interested
- 155 in running M, it is important to make this process transparent and robust. The representation of M as a computer program typically means that in the set of program variables  $\mathcal{V} = \mathcal{A} \cup \mathcal{P}$  there is a subset  $\mathcal{P}$  of *passive* variables that do not carry derivative information, such as variables of non-differentiable type (integers, strings) but also floating-point parameters, physical constants, or descriptors of the problem domain. The size of the trace T is a major factor on the computational
- 160 efficiency of the adjoint. Thus, one must categorize program variables into  $\mathcal{A}$  and  $\mathcal{P}$  with the aim to minimize the set  $\mathcal{A}$  of active (i.e. type-changed) program variables. In ISSM this is accomplished by changing double variables to a type named for brevity e.g. TA (in ISSM IssmDouble) for variables in  $\mathcal{A}$  and to e.g. TP (in ISSM IssmPDouble) for variables in  $\mathcal{P}$ , respectively. This categorization must be performed by an expert familiar with the global data dependencies in  $\mathcal{M}$
- 165 and the work that has been done for Mand. In addition, it can be incrementally done for code contributions supplied as double variables without breaking M. This permits easy tracking as is a key aspect of the overloaded AD approach that is amenable to easy maintenance and expansion of a complex C++ framework such as ISSM, and one of the key reasons why this approach was selected to compute M in ISSM. In the end, all double types should eventually be replaced with
- 170 either TA or TP. Following good practices in C++ both TA and TP are typedef ed in a central location switched only there via a preprocessing macro to use the active type supplied by the AD tool or hide the distinction altogether for running the plain M. This implies that one cannot explicitly overload methods or define data structures for distinct TA and TP without filtering the declaration and definition of the TA variant by the preprocessor as well. Such code duplication introduces an
- 175 undue code maintenance burden.

Instead, the recommended approach is the use of C++ template classes and methods where TA and TP are the concrete arguments for an abstract template type T. While the use of templates may imply a larger up-front effort if they weren't present in the model code before, as was the case in ISSM, the long term benefits not only for  $\overline{M}$  but also the plain M are obvious when one wants to

180 experiment with computation at different precision levels. Aside from ISSM, an established general example for this kind of templating is the use of SACADO within Trilinos (The Trilinos Project; Phipps et al., 2008). Of particular value in ISSM was the migration of data containers for arrays and (sparse) matrices and that of existing wrappers to malloc and free to templated wrappers

of new and delete, in ISSM called xNew and xDelete. The latter is necessary not only to

185 properly instantiate TA variables but also permits specializations for TA in  $\overline{M}$  as explained in the following.

Overall, the use of C++ template classes introduces an additional element of complexity to the code, but it also means that the same exact framework can be used to compute both the forward run and its gradient, using the exact same code base, without the need to manage different code for the

190 computation of  $\overline{M}$ . This in turn means that ISSM can be compiled using ADOLC and shipped to the community in only one binary version that can accomodate both M and  $\overline{M}$ , which is a critical feature in terms of extending AD features to the wider community without having to dive into the intrinsic complexities of AD transformation itself.

Interpreting the trace by R(T) means that the space holding the data for the variables  $r, a, b \in \mathcal{V}^*$ 

- 195 (the set of instantiated program variables at run time) occurring in each operation  $r = \phi(a, b)$  must be represented by some mapping  $\omega : \mathcal{V}^* \mapsto \mathbb{N}^+$  to pseudo addresses. In ADOL-C these addresses are called *locations* and represent indices in a work array held by R(T). The pseudo addresses must be managed through the TA constructor and destructor in a fashion similar to the memory management of the actual program variables themselves, i.e. pseudo addresses are assigned from and returned
- 200 to a pool  $\Omega$  of available addresses. However, no distinction between heap and stack variables is made and generally data locality will not be preserved. On the other hand, for special operations  $\phi_A$ with array arguments of size *s*, that is, calls to external solvers (see section 2.3) or MPI routines (see section 2.4), it would be counterproductive to record in *T* the pseudo-addresses of each of the *s* array elements rather than a consecutive range. The latter, however, imposes that the pool be primed by
- 205 some call  $\Omega(s)$  such that  $\omega$  returns consecutive pseudo addresses when called by the constructor for each element in the array. In ADOL-C this is done by calling ensureContiguousLocations immediately before an active array is instantiated. Avoiding copying of array values in M as an efficiency measure means that any given array has a good chance of being used in  $\phi_A$ , and to avoid littering the code with preprocessor-guarded calls to  $\Omega(s)$ , we decided in ISSM to instead add a call
- 210 to ensureContiguousLocations in the TA specialization of xNew:

```
#if defined(_HAVE_ADOL_C_)
template <>
adouble* xNew(unsigned int s) {
    ensureContiguousLocations(s);
215 adouble* aT_p=new adouble[s];
    assert(aT_p);
    return aT_p;
    }
    #endif
```

Thus, the use of  $\phi_A$  means frequent calls to  $\Omega(s)$ . The implementation of  $\Omega(s)$  searches the pool to find a sufficiently large consecutive range or otherwise grow the pool by at least *s* addresses. In ISSM, this search became a significant performance bottleneck. Minimizing the search effort and balancing it with pool growth has been newly implemented in ADOL-C and is controlled by setting parameters via setStoreManagerControl for the ratio of the addresses inside and outside of the pool and the maximal pool size to trigger searches.

#### 2.3 Binding to External Solvers

The strength of the AD approach lies in the automated differentiation of all the parts of the model that are changing and expanding as the model is being developed as opposed to manually programming adjoints which is error prone and implies a significant effort in terms of code development and maintenance. However, because models reference libraries, for *fixed* mathematical mappings, that may have easily programmable high-level adjoints, one can and should exploit this insight as it can yield efficiency gains not obtainable by an AD tool remaining unaware that calls to some

collection of subroutines represent a certain mathematical mapping. A good example are linear

- solver libraries that may not even be written in the same programming language as the model in 235 question. Considering the system As = rAs = b, the solver S computing  $s = S(A, r) \cdot s = S(A, b)$ needs an adjoint counterpart solving  $A^T t = \bar{s}$ , incrementing  $\bar{r} = \bar{r} + t \cdot \bar{b} = \bar{b} + t$  and, if  $A \in A$ , also  $\bar{A} = \bar{A} - \bar{b}s^T$ . Clearly, the solve with the transposed can easily be implemented reusing the solver S as in  $t = S(A^T, \bar{s})$ . However, the answers to the following questions affect the efficiency and ought to be considered in the context of the given model:
- 240 Q1: Is  $A \in A$  which therefore requires the rank-1 update to  $\overline{A}$ ?

245

250

- Q2: Are the previous values of s, r and A overwritten by S but used to compute partial derivatives for Eq. (3) and therefore must be saved and restored?
- Q3: If S is a direct solver should one save the factors or refactorize for the adjoint solve?
- Q4: How does the external function interface used by R(T) allow for efficient reuse of intermediate buffers?

External functions  $f_e$  that had been supported by ADOL-C had the signature  $f_e(l_x, x, l_y, y)$  with inputs x, outputs y for the original call,  $l_x, l_y$  their respective array lengths, and  $\overline{f_e}(l_y, \overline{y}, l_x, \overline{x})$  for the adjoint counterpart. In the case of a linear system with  $A \in A$ , the input x is packed with both A and r while y contains the solution s of the system on return. This, however, was insufficient for binding to solvers from the GNU Scientific Library (GSL) (Galassi, 2009) and MUMPS (Amestoy et al., 2001) used by ISSM. The following extensions were implemented in ADOL-C to enable the

E1: Expanded the  $\overline{f_e}$  interface by x, y to  $\overline{f_e}(l_y, \overline{y}, l_x, \overline{x}, x, y)$  to enable refactoring;

ISSM adjoints but are generic in nature and would need to be supported in some form by other tools.

- E2: added optional parameters to pass in the sparsity pattern of A for MUMPS as a generic integer array i of length  $l_i$  for both  $f_e$  and  $\overline{f_e}$ ;
- E3: added controls for storing and restoring prior values of x and y;
- E4: added tracking of the maximal  $l_x, l_y$  in sequences of  $f_e$  calls.

Regarding Q1 we know that in ISSM  $A \in A$ , and regarding Q2 the parameters passed to  $f_e$  have no other uses and therefore, using the controls (E3), we avoid (re)storing their values. The direct 260 solver from the GSL used here had no API control to back solve with the factors for the transposed and we did not want to reverse engineer the permutation representation. Hence the refactoring was done as a matter of convenience for the sequential reference case requiring E1. The parallel and therefore practically more efficient MUMPS solver operates on sparse, distributed A, therefore requiring E2. MUMPS offers both the ability to store the factors to file and perform the back-solve for

- the transpose. However, the MUMPS portion of the runtime is comparatively small (see section 3). 265 Consequently the overhead for the file I/O when considering the factor data size after fill-in is not expected to yield much practical benefit in this context, answering Q3. Finally, for Q4, the extension E4 is exploited because transient runs of ISSM need to account for changes in the system size and a preallocation with the maximal buffer sizes therefore avoids some of the memory management 270 overhead.

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#### Handling Parallelism with the AdjoinableMPI Library 2.4

As is the case with ISSM, practically relevant science problems incur a computational complexity that necessitates execution on parallel hardware, often using MPI as the paradigm of choice. Sending data with MPI from a source buffer s to a target buffer t can be interpreted as a simple assignment

- t = s. This implies for the adjoint an increment  $\bar{s} = \bar{s} + \bar{t}$ , that is, the adjoint is a communication of a 275 communication is a reversal of the data flow between buffers. Calculating the adjoint of the adjoint values in the reverse direction t = s means sending back the value of  $\bar{t}$  and computing  $\bar{s} = \bar{s} + \bar{t}$ . The principles for adjoining two-sided MPI communication have been explored in Utke et al. (2009). The development of the AdjoinableMPI (AMPI) wrapper library AdjoinableMPI started in Fall of 2012.
- It is designed to provide an AD tool-independent implementation for adjoining MPI-parallelized C. 280 C++ and Fortran models. The wrapper interfaces distinguish themselves from the original MPI by the prefix AMPI and have a few additional parameters where needed to enable the adjoint functionality. AMPI also provides additional types and predefined symbols. Discussing the internal design of AMPI is outside the scope of this paper. However, the application to ISSM is the first large scale
- 285 practical use of AMPI and in the following we will discuss the steps taken to use it in the ISSM code base.

For testing and small scale experiments, the ISSM code, as many other models, treats MPI parallelization as a compile-time option controlled by preprocessor macros. Furthermore, turning the

Table 1. Logic variants encapsulated in the ISSM MPI wrapper library

	MPI	AD	
1	no	no	emulate MPI semantic, e.g. with memcpy between buffers for MPI_Reduce
2	no	yes	emulate MPI semantic but for differentiation need to use the overloaded assignments
			instead of memcpy
3	yes	no	pass-through to plain MPI
4	yes	yes	call AMPI routines

290

adjoint capability on and off as suggested in section 2.2 would imply additional switching between the original MPI and AMPI with code duplication and the potential for errors. To avoid these undesirable consequences we decided to introduce in ISSM another wrapper layer (prefixed ISSM\_MPI ) of calls and definitions to encapsulate completely the four functionality variants listed in Table 1. The approach is shown for MPI\_Reduce in Fig. 1. Variant 1 is implemented by line 21, variant 2

```
1 int ISSM_MPI_Reduce(void *sendbuf, void *recvbuf, int count,
2
                ISSM_MPI_Datatype datatype, ISSM_MPI_Op op, int root,
3
                ISSM_MPI_Comm comm) {
4
          int rc=0;
5
  #ifdef _HAVE_MPI_
6
   #ifdef _HAVE_AMPI_
7
          rc=AMPI_Reduce(sendbuf, recvbuf, count, datatype, op, root, comm);
8
   #else
9
          rc=MPI_Reduce(sendbuf, recvbuf, count, datatype, op, root, comm);
10 #endif
11 #else
12 #ifdef _HAVE_ADOLC_
13
          if (datatype==ISSM_MPI_DOUBLE) {
14
             IssmDouble* activeSendBuf=(IssmDouble*)sendbuf;
15
             IssmDouble* activeRecvBuf=(IssmDouble*)recvbuf;
16
             for(int i=0;i<count;++i) activeRecvBuf[i]=activeSendBuf[i];
18
          }
19
          else
20 #endif
21
               memcpy(recvbuf, sendbuf, sizeHelper(datatype)*count);
22 #endif
23
      return rc;
24 }
```

Figure 1. Code for wrapping reduction

by lines 13-18 and 21, variant 3 by line 9, and variant 4 by line 7. The example code reflects some of the simplifying assumptions being made based on the MPI usage patterns in ISSM. Examples are the

	ISSM_MPI_DOUBLE #define to	IssmDouble (=TA) typedef to	ISSM_MPI_PDOUBLE #define to	IssmPDouble (=TP) typedef to
1	2	double	3	double
2	2	adouble $\frac{*+}{\sim}$	3	double
3	MPI_DOUBLE	double	MPI_DOUBLE	double
4	$\operatorname{AMPI}_{\operatorname{ADOUBLE}} \stackrel{**++}{\sim}$	adouble $\frac{*}{\sim}$	MPI_DOUBLE	double

**Table 2.** Types per variant from Table 1;  $^{++}$  active type from ADOL-C,  $^{+++}$  active MPI type from AMPI

absence of MPI\_IN\_PLACE and user-defined MPI types as well as all active data being of double precision. This permits the simple distinction on line 12 but of course the logic covers all other non-differentiated MPI types that may occur. The following lines 14,15 indicate the pairing between the C++ type definition IssmDouble and the wrapper-defined ISSM\_MPI\_DOUBLE to signify active data being communicated when AD is enabled. Mirroring in MPI the approach described in section 2.2 with two types TA and TP, the passive type TP (concrete IssmPDouble) has an MPI counterpart defined in the wrapper as ISSM\_MPI\_PDOUBLE. It too can be the value passed in via the datatype argument at line 1 and signals to AMPI passive floating point communications. The

305 Table 2.

300

The matching between the actual type of the buffer and the corresponding MPI type must extend to the templating approach suggested for the type change in section 2.2. The main reason is to avoid type errors, which become more common as the complexity of a code increases. Because the MPI standard keeps the definition of MPI Datatype opague and MPI types can be created at runtime,

definitions provided by the ISSM\_MPI wrapper corresponding to types TA and TP are given in

- 310 that is, the value the value of MPI\_Datatype cannot be used as the value template parameter itself (as it may not be a compile time constantusable as a value template parameter). Therefore, a template buffer type T must be paired with the corresponding MPI type value declared as a pointer parameter. The approach is shown, as show in Fig. 2. Assuming many AMPI calls in foo, this reduces the count of confines the code locations where type errors may be introduced to the template
- 315 instantiations themselves.

Finally, a practical concern for using the parallelized adjoint is the handling of sensitivities to quantities that are uniformly initialized across ranks, such as parameters. Frequently, as was the case within ISSM, these quantities are initialized from files or otherwise per process in the parallel case the same way as in the sequential case. In the parallel case that implies a replication of the

same quantity across ranks. However, to obtain the correct sensitivities, the quantity q in question, should be unique, in other words that quantity must be uniquely initialized at one root rank and then broadcast to the other ranks. Otherwise, for r ranks, then at each rank one would obtain only a part  $\bar{q}_i$  of the total  $\bar{q}$  and would have to "manually" sum up  $\bar{q} = \sum \bar{q}_i$ . With an initial broadcast of q,

```
// this declaration comes from the wrapper:
extern ISSM_MPI_Datatype ourISSM_MPI_DOUBLEVal,ourISSM_MPI_DOUBLEVal;
template <class T, ISSM_MPI_Datatype *mt_p> void foo(T *t){
    ISSM_MPI_Reduce(t, ..., ,*mt_p,..);
}
void bar(IssmDouble *x, IssmPDouble *p) {
    foo<IssmDouble,&ourISSM_MPI_DOUBLEVal>(x);
    foo<IssmPDouble,&ourISSM_MPI_PDOUBLEVal>(p);
}
```

Figure 2. Code snippets for templating with corresponding MPI data types

however, the corresponding adjoint provided through AMPI by using AMPI\_Bcast is that exact sum reduction and R(T) yields the correct adjoint at the broadcast root. This notion similarly applies to any situation where a conceptually unique quantity of active type is implicitly replicated on some ranks.

#### 2.5 Validation

ISSM is validated in AD mode by continuously running a test suite within the Jenkins (Smart, 2011)
integration and development framework (available here at http://issm.jpl.nasa.gov/developers/). A detailed description of the suite of benchmarks is given in Table 3) The aim is to 1) compare forward runs in ISSM with their counterparts when overloaded operators are switched on. The results should be identical within double precision tolerances; and 2) compare forward and reverse runs carried out with ISSM AD on and off, using the GSL and MUMPS solvers. Comparisons of gradients
computed in AD mode with standard forward differences methods are also carried out to make sure the gradients computed (essentially in reverse scalar mode, which is the mode of predilection in ISSM for data assimilation) in AD mode are accurate.

#### **3** Application

#### 3.1 Application

- 340 The ongoing use of adjoint computations includes sensitivity studies and state estimation problems for transient model runs. Because this paper concentrates on technical aspects, we show, only as exemplary evidence of the practical usefulness, some sensitivities of cost functions calibrated for two sensors commonly encountered in Cryosphere Science, altimeters (that measure surface elevation), and radars (that measure surface displacement, or velocity). In Fig. 3, the sensitivity (gradient) of a
- 345 cost function related to ICESat-1 altimetry with respect to Surface Mass Balance (SMB) is shown

**Table 3.** ISSM AD validation suite integrated within Jenkins for continuous integration and delivery (Smart, 2011). Tests 3001 to 3010 and 3101 to 3110 test the repeatability of forward runs with and without ADOL-C compiled, but with no AD drivers specifically called. The forward runs involved are the standard stress balance, mass transport and thermal solutions, with 2D SSA (Shelfy-Stream Approximation MacAyeal (1989)), 3D SSA, 3D HO (Higher-Order, (Blatter, 1995; Pattyn, 1996)), 3D Full-Stokes (Stokes, 1845) and DG (Discontinuous-Galerkin) formulations. Test 3015 tests the AD GSL capability, by comparing the AD Forward Scalar mode (where we compute the gradient of ice volume with respect to ice thickness at one vertex of the mesh) against a simple forward differences computation on the same gradient. Test 3019 validates the AD GSL capability in Reverse Scalar mode (where we compute ice volume gradient with respect to thickness at all vertices of the mesh) vs the Forward Vectorial mode. Both gradients should be identical. Finally, test3119 validates the parallel capabilities (AD MUMPS) by comparing reverse scalar computations with GSL and MUMPS.

Test	Solution Sequence	Formulation	Solver	Description
3001/3101	Stress Balance	2D SSA	GSL/MUMPS	Equal runs with overload on and off
3002/3102	Stress Balance	3D SSA	GSL/MUMPS	Equal runs with overload on and off
3003/3103	Stress Balance	3D HO	GSL/MUMPS	Equal runs with overload on and off
3004/3104	Stress Balance	3D FS	GSL/MUMPS	Equal runs with overload on and off
3005/3105	Mass Transport	2D	GSL/MUMPS	Equal runs with overload on and off
3006/3106	Mass Transport	2D (DG)	GSL/MUMPS	Equal runs with overload on and off
3007/3107	Mass Transport	3D	GSL/MUMPS	Equal runs with overload on and off
3008/3108	Thermal Steady State	3D	GSL/MUMPS	Equal runs with overload on and off
3009/3109	Thermal Transient	3D	GSL/MUMPS	Equal runs with overload on and off
3010/3110	Transient 2D	2D	GSL/MUMPS	Equal runs with overload on and off
3015	Mass Transport 2D	2D	GSL	AD Forward Scalar vs Forward Differences
3019	Thermal Transient	3D	GSL	AD Reverse Scalar vs AD Forward Vectorial
3119	Thermal Transient	3D	MUMPS vs GSL	AD Reverse Scalar (MUMPS vs GSL)Scalar

for three epochs, September 2003, June 2006 and February 2009. The cost function <u>J<sub>Altimetry</sub></u> is the spatio-temporal average of the misfit between modeled surface elevations (from the mass-transport module of the transient solution in ISSM), and the corresponding ICESat-1 altimetry record. This cost function was used within an inversion method to reconstruct SMB over the entire ICESat-1 time

- 350 record (Larour et al., 2014). This example would correspond to the case where function f from Eq. 1 is a transient ice-flow model (including mass transport and stress-balance, no thermal, described in Larour et al. (2014)), fully non-linear in its treatment of the material rheology, and highly resolved both spatially (km resolution at the coastline) and temporally (two-week time step). The final output (corresponding to y in Eq. 1) of the model is the altimetry cost function  $J_{Altimetry}$  itself and the
- 355 model input SMB (corresponding to x in Eq. 1). The sensitivity displayed in Fig.3 is therefore the AD computed  $\frac{\partial J_{Altimetry}}{\partial SMB}$  at different time steps corresponding to the transient model input SMB.



**Figure 3.** Gradient of surface altimetry cost function (spatio temporal average of the misfit between the 6 year record of observed ICESat altimetry and modeled surface elevation) with respect to SMB (in years). The gradient is computed for every time step of a two-week interval time series between 2003 and 2009. We only show three periods (September 2003, June 2006 and February 2009) in the entire interval. The location for the study is North-East Greenland Ice Stream, and the gradient is taken from an inversion study of the surface forcings necessary to best fit the ICESat altimetry record (Larour et al., 2014).

A second type of observations garnering considerable interest is temporally resolved time series of radar observations (using speckle-tracking, or InSAR to infer surface deformation of the ice) to measure variations of surface velocity on a seasonal time scale (one observation every 2 months, as

- in Moon et al. (2015)). Temporally inverting for such time series, trying to reconstruct for example basal friction is a topic of interest due to its relevance in terms of understanding the dynamics of calving, basal slip and shear softening, and associated feedback mechanisms. In Fig. 4, we demonstrate the feasibility of such inversions by computing the gradient of a cost function  $J_{Velocity}$  related to the Moon et al. (2015) time series with respect to basal drag  $\alpha$  at epochs 2008, 2010 and 2013.
- 365 The cost function is the spatio-temporal average of the misfit between modeled surface velocity (from the stress-balance module of the transient solution in ISSM) and observed surface velocities. This example would correspond to the case where function f from Eq. 1 is the same transient ice-flow model described previously, the final output (corresponding to y in Eq. 1) of the model is the velocity cost function  $J_{Velocity}$  itself and the model input basal friction  $\alpha$  (corresponding to x in
- 370 Eq. 1). The sensitivity displayed in Fig.4 is therefore the AD computed  $\frac{\partial J_{Velocity}}{\partial \alpha}$  at different time steps corresponding to the transient model input  $\alpha$ .



**Figure 4.** Gradient of surface velocity cost function (spatio temporal average of the misfit between a 5 year time series of observed surface velocities and the modelled surface velocity) with respect to the basal drag coefficient ( in  $yr^{1/2}/m^{5/2}$ ). The gradient is computed for every time step of a two-week interval time series between 2008 and 2013. We only show three periods (2008, 2010 and 2013) in the entire interval. The location for the study is Upernavik Glacier, Central West Greenland, and the gradient is taken from an inversion study of the basal forcings required to match observed RADAR time series of surface velocites and calving front positions (*Larour et al, 2016, in preparation*).

#### 3.2 Benchmarking

Currently, these adjoint computations are performed on the Pleiades cluster at the NASA Advanced Supercomputing Center. Given the theoretical invariance of fact that the effort for computing the

- adjoint  $\nabla f$  with respect to the gradient size *n* as stated in section 2 relies on one reverse sweep, the practical question remains about becomes what the actual runtime overhead of computing both *f* and  $\nabla f$  when compared to computing just the original *f* is. Here in particular one considers the effects of compiler optimization on the original code for *f* using the built-in floating point operations on one side and in contrast to that the overhead incurred by calling the overloaded operators as well as
- the creation, storage, and interpretation of the trace T (as the means to compute  $\nabla f$ ) on the other. Aside from the fact that the tuning of the ISSM adjoint is a work in progress we want to highlight the overwhelming impact of the application specific aspects on the runtime ratio. Therefore, we emphasize that exhibiting any particular overhead number as the ultimate result of scenario-specific tuning is of little use to the practitioner wanting to answer science questions. Rather, using examples from the ISSM work we want to show what may prevent achieving a satisfactory overhead.

The earliest ISSM adjoint computations took place before the MPI wrapper library was introduced and therefore were done sequentially with the GSL solvers. To evaluate the performance, a representative test case was chosen from the ISSM regression test suite<del>(test 101, modeling the stress-balance</del>, test 101. This test models the steady state stress balance of a square ice shelf <del>). The of size 50km</del>

390 x 50km and thickness 1000 m (at the grounding line) to 300 m (at the ice front). The ice shelf is

clamped at the grounding line, free to flow at the ice front (one one side only), with non-linear viscosity dependent on Glen's law Glen (1955, 1958). The formulation relies on the Shelfy-Stream Approximation (SSA, MacAyeal (1989)), and is inherently non-linear, relying on a Picard iteration for convergence. For more details on this test, we refer the reader to Larour et al. (2012). For

- 395 <u>benchmarking, the</u> problem size was indirectly set by specifying a measure for the maximal resolution of the generated mesh, thus increasing the number of mesh elements generated for a smaller resolution parameter. The <u>parallel runs were carried out using 8 cpus</u>. The overhead factors are shown separately for the overloading as such and the generation of T and R(T) as wall-clock comparisons to the unmodified model compiled with default optimization (-O2) in Fig. 5 (upper frame). While
- 400 this plot indicates a small overhead factor of  $\approx 4.5$  in particular for the largest mesh case (distance 12.5 km) the reason for this becomes apparent in the plot on the lower frame. It shows that the majority of the run time is consumed by the GSL solver (libgsl) completely overshadowing any of the overhead caused by the adjoint for the largest mesh. We want to emphasize that GSL was chosen not for its efficiency but for the simplicity of the setup which quickly enabled adjoint computations.
- 405 Introducing AMPI and thereby moving to a more appropriate solver (MUMPS) causes the adjoint overhead to become more prominent. The most consequential change necessitated by the use of MPI is the forced contiguity of the pseudo addresses (see section 2.2).



**Figure 5.** Test problem setup with sequential GSL solver; plots over max mesh distance show overhead factors (upper frame), and approximate run-time portions of the model execution with adjoint computation (lower frame)

In combination with the much reduced impact of the solver, the overhead factor for equivalent test problems reached temporarily up to 145 which clearly was not acceptable. Subsequent changes

- 410 to the ADOL-C tool included improvements in the internal address management. Changes to the model included modifications of the sparse data format, enabling the control of the I/O buffering for the trace T and the ADOL-C address manager via the ISSM configuration specific to the setup to be computed. The combination of these changes led to regaining better performance and effective overall overheads between 10 and 30. Analyzing the details of the performance shows that the overhead
- 415 factor for the trace creation and interpretation does not change significantly (see Fig. 6 bottom) with the mesh size in accordance with the theoretical result for the adjoint computation. The large majority of the total overhead, evidenced by the runtime portion for libadolc in Fig. 6 (top) still originates in the internal address management. While the overall overhead factors are sufficiently small for the practical use of the adjoints for science problems further improvements in the addressing scheme are



420 clearly warranted and subject of ongoing work.

**Figure 6.** Test problem setup with parallel MUMPS solver; plots over max mesh distance show approximate run-time portions of the model execution with adjoint computation (upper frame) and overhead factors for the trace creation and interpretation (lower frame)

#### 4 Conclusions

We developed a new adjoint capability in ISSM, based on the ADOL-C framework and the AdjoinableMPI library, which is to our knowledge the first time this type of approach has been systematically applied to a software framework of this size and complexity. Despite the difficulties

- 425 encountered rewriting the software, the overloaded approach is transparent to the user, which is critial critical given the size of the larger Cryosphere Science community community that is not familiar with the adjoint work, and for which classic approaches such as source-to-source transformation have proven to be overly may not be applicable, or may prove prove more cumbersome. The flexibility of this approach allows in particular for quick turn-around in developing adjoint models
- 430 of new parameterizations which are not easily hand derived. This is a major advantage in that it opens this approach to the wider community. This, given the large amount of remote sensing data currently being collected and under-utilized, could prove paramount if we are to hindcast validate projections of sea-level rise. Further work is of course required to bring in additional observations such as gravity sensors, or radar stratigraphy observations, which will involve development of new
- 435 cost functions, and scalability in 3D. Though this is complex in that it requires integrated resiliency and adjoint checkpointing schemes for long running transient modeling scenarios, our approach has proven flexible, and should lead to a brand new set of data assimilation capabilities that have already been available to other Earth Science communities for a long time. Indeed, by allowing temporal data assimilation for a large number of sensors and models, such as demonstrated here with the use 440 of altimetry and radar sensors for mass transport and stress balance models respectively, ISSM paves
- the way for wider integration between the modeling and observational Cryosphere community.

#### 5 Code Availability

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The ISSM code and its AD components are available at http://issm.jpl.nasa.gov. The instructions for the compilation of ISSM in AD mode, along with test cases is presented in the supplement attached to this manuscript.

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