

POM.gpu-v1.0: a GPU-based Princeton Ocean Model

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Abstract. Graphics Processing Units (GPUs) is an attractive solution in many scientific applications due to its high performance. However, most existing GPU conversions of climate models use GPUs for only a few computationally intensive regions. In the present study, we redesign the mpiPOM (a parallel version of the Princeton Ocean Model) with GPUs. Specifically, we first convert the model
5 from its original Fortran form to a new Compute Unified Device Architecture C (CUDA-C) code, then we optimise the code on each of the GPUs, the communications between the GPUs, and the I/O between the GPUs and the Central Processing Units (CPUs). We show that the performance of the new model on a workstation containing 4 GPUs is comparable to that on a powerful cluster with 408 standard CPU cores, and it reduces the energy consumption by a factor of 6.8.

10 1 Introduction

High-resolution atmospheric, oceanic and climate modelling remain significant scientific and engineering challenges because of the enormous computing, communication, and storage requirements involved. Due to the rapid development of computer architecture, in particular the development of multi-core and many-core hardware, the computing power that can be applied to scientific problems
15 has increased exponentially in recent decades. Parallel computing methods, such as the Message Passing Interface (MPI, Gropp et al. (1999)) and Open Multi-Processing (OpenMP, Chapman et al. (2008)) have been widely used to support the parallelization of climate models. However, supercomputers are becoming increasingly heterogeneous involving devices such as the GPU and the Intel Many Integrated Core (Intel MIC), and new approaches are required to effectively utilize the new
20 hardware.

In recent years, a number of scientific codes have been ported to the GPU as shown in Table 1. Most existing GPU acceleration codes for climate models are only operating on certain hot spots of the program, leaving a significant portion of the program still running on CPUs. The speed of some

Table 1. Existing GPU porting work in climate fields. The speedups are normalized to one CPU core.

Model Name	Model Description	Porting Modules to GPU	Speedup
WRF	Weather research and forecasting	WSM5 microphysics	8
WRF-Chem	WRF chemical	Chemical kinetics kernel	8.5
POP	Parallel ocean program	Loop structures	2.2
COSMO	Consortium for small-scale modelling	Dynamical core	22.7
NIM	Non-hydrostatic icosahedral model	Dynamical core	34
ASUCA	Non-hydrostatic weather model	Dynamical core & physical	80

subroutines reported in the Weather Research and Forecast (WRF) (Michalakes and Vachharajani,
25 2008) and WRF-Chem(Linford et al., 2009) is improved by a factor of approximately 8 whereas
the whole model achieves limited speedup because of partial porting. The speed of POP (Zhenya
et al., 2010) is improved by a factor of only 2.2 because the model only accelerated a number of loop
structures using the OpenACC Application Programming Interface (OpenACC API). The speed of
COSMO(Leutwyler et al., 2014), NIM(Govett et al., 2010) and ASUCA(Shimokawabe et al., 2010)
30 are greatly improved by multiple GPUs. We believe that the elaborate optimization of the memory
access of each GPU and the communication between GPUs can further accelerate these models.

The objective of our study was to shorten the computation time of the Princeton Ocean Model
(POM) by parallelizing its existing model structures using the GPU. Taking the parallel version of
the Princeton Ocean Model (mpiPOM), we demonstrate how to code an ocean model so that it runs
35 efficiently on GPU architecture. We first convert the mpiPOM from its original Fortran version into
a new Compute Unified Device Architecture C (CUDA-C) version, POM.gpu-v1.0. CUDA-C is the
dominant programming language for GPUs. We then optimise the code on each of the GPUs, the
communications between the GPUs, and the I/O between GPUs and the CPUs to further improve
the performance of POM.gpu.

40 To understand the accuracy, performance and scalability of the POM.gpu code, we customized
a workstation with four Nvidia K20X GPUs. The results show that the performance of POM.gpu
running on this workstation is comparable to that on a powerful cluster with 408 standard CPU
cores.

This paper is organized as follows. In Section 2, we review the mpiPOM model. In Section 3,
45 we briefly introduce the GPU computing model. In Section 4, we present the detailed optimization
techniques. In Section 5, we report on the correctness, performance and scalability of the model. We
present the code availability in Section 6 and conclude our work in Section 7.

2 The mpiPOM

The mpiPOM is a parallel version of the POM. It retains most of the physics of the original POM (Blumberg and Mellor, 1983, 1987; Oey et al., 1985a, b, c; Oey and Chen, 1992a, b), and includes satellite and drifter assimilation schemes from the Princeton Regional Ocean Forecast System (Oey, 2005; Lin et al., 2006; Yin and Oey, 2007), Stokes drift and wave-enhanced mixing (Oey et al., 2013; Xu et al., 2013; Xu and Oey, 2014). The POM code was reorganized and the parallel MPI version was implemented by Jordi and Wang (2012) using a two-dimensional data decomposition of the horizontal domain. The MPI is a standard library for message passing and it is widely used to develop parallel programs. The POM is a powerful ocean model that has been used in a wide range of applications: circulation and mixing processes in rivers, estuaries, shelves, slopes, lakes, semi-enclosed seas and open and global oceans. It is also at the core of various real-time ocean and hurricane forecasting systems, e.g., the Japanese coastal ocean and Kuroshio current (Miyazawa et al., 2009; Isobe et al., 2012; Varlamov et al., 2015); the Adriatic Sea Forecasting System (Zavatarelli and Pinardi, 2003); the Mediterranean Sea forecasting system (Korres et al., 2007); the GFDL Hurricane Prediction system (Kurihara et al., 1995, 1998), the US' Hurricane Forecasting System (Gopalakrishnan et al., 2010, 2011) and the Advanced Taiwan Ocean Prediction system (Oey et al., 2013). Additionally, the model has been used to study various geophysical fluid dynamical processes (e.g., Allen and Newberger, 1996; Newberger and Allen, 2007a, b; Kagimoto and Yamagata, 1997; Guo et al., 2006; Oey et al., 2003; Zavatarelli and Mellor, 1995; Ezer and Mellor, 1992; Oey, 2005; Xu and Oey, 2011, 2014, 2015; Chang and Oey, 2014; Huang and Oey, 2015; Sun et al., 2014, 2015). For a more complete list, please visit the POM website (<http://www.ccpo.odu.edu/POMWEB>).

The mpiPOM experiment used in this paper is one of two that were designed and tested by Professor Oey and students; the codes and results are freely available at the FTP site (<ftp://profs.princeton.edu/leo/mpipom/atop/tests/>). The reader can refer to Chapter 3 of the lecture notes (Oey, 2014) for more detail. The test case is a dam-break problem in which warm and cold waters are initially separated in the middle of a zonally periodic channel (200km \times 50km \times 50m) on an f-plane, with walls at the northern and southern boundaries. Geostrophic adjustment then ensues and baroclinic instability waves amplify and develop into finite-amplitude eddies in 10~20 days. The horizontal grid sizes are 1 km and there are 50 vertical sigma levels. Although the problem is a test case, the code is the full mpiPOM version used in the ATOP forecasting system.

The model solves the primitive equation under hydrostatic and Boussinesq approximations. In the horizontal, spatial derivatives are computed either using centred-space differencing or Smolarkiewicz's positive definite advection transport algorithm (Smolarkiewicz, 1984) on a staggered Arakawa C-grid; both schemes have been tested, but the latter is reported here. In the vertical, the mpiPOM supports terrain-following sigma coordinates and a fourth-order scheme option to reduce the internal pressure-gradient errors (Berntsen and Oey, 2010). The mpiPOM uses the time-splitting technique to separate the vertically integrated equations (external mode) from the vertical structure

85 equations (internal mode). The external mode calculation is responsible for updating the surface elevation and vertically averaged velocities. The internal mode calculation updates the velocity, temperature and salinity, as well as the turbulence quantities. The three-dimensional internal mode and the two-dimensional external mode are both integrated explicitly using a second-order leapfrog scheme. These two modules are the most computationally intensive kernels of the mpiPOM model.

90 The main computational problem of the mpiPOM is memory bandwidth limited. To confirm this issue, we use the runtime performance API tool to estimate the floating point operation count and the memory access instruction count, as in Browne et al. (2000). The results reveal that the computational intensity, defined as floating point operations per byte transferred to or from memory, of the mpiPOM is approximately 1:3.3, whereas the computational intensity provided by a modern high-
95 performance CPU (an Intel SandyBridge E5-2670) is 7.5:1. Many large arrays are mostly pulled from main memory and there is poor data reuse in the mpiPOM. In addition, there are no obvious hot spot functions in the mpiPOM, and even the most time-consuming subroutine occupies only 20% of the total execution time. Therefore, porting a handful of subroutines to the GPU is not helpful in improving the model efficiency. This explains why we must port the entire program from the CPU
100 to the GPU.

3 GPU computing model overview

Modern GPUs employ a stream-processing model with parallelism. Each GPU contains a number of stream multiprocessors (SMs). In this work, we carried out the conversion using four Nvidia's K20X GPUs. Each K20X GPU contains 14 SMs and each SM has 192 single precision processors
105 and 64 additional processors for double precision. Although the computational capability of each processor is low, one GPU with thousands of processors can greatly boost the performance compared to the CPU. In computing, FLOPS (FLoating-point Operations Per Second) is a measure of computer performance. The theoretical peak performance of each K20X GPU is 3.93 teraFLOPS (TFLOPS, one trillion floating-point operations per second) for the single precision floating-point calculations.
110 In contrast, a single Intel SandyBridge E5-2670 CPU is only capable of 0.384 TFLOPS.

Each pair of GPUs shares 6 Gigabytes (GB) of memory, with the interface having a potential bandwidth of 250 GB/s. Figure 3 illustrates the memory hierarchy of the K20X GPU. Each SM possesses some types of fast on-chip memory such as register, L1 cache, shared memory and read-only data cache. In GPUs, the register is the fastest memory, of which the size is 256 Kilobytes
115 (KB) for each SM. The shared memory and the L1 cache use the common 64 KB space which can be partitioned as 16/48 KB, 32/32 KB or 48/16 KB. The 48 KB read-only data cache is useful for holding frequently used values that remain unchanged during each stage of the processing.

There are three widely used methods for porting a program to GPUs. The first method uses drop-in libraries provided by CUDA to replace the existing code, as in Siewertsen et al. (2012). The second

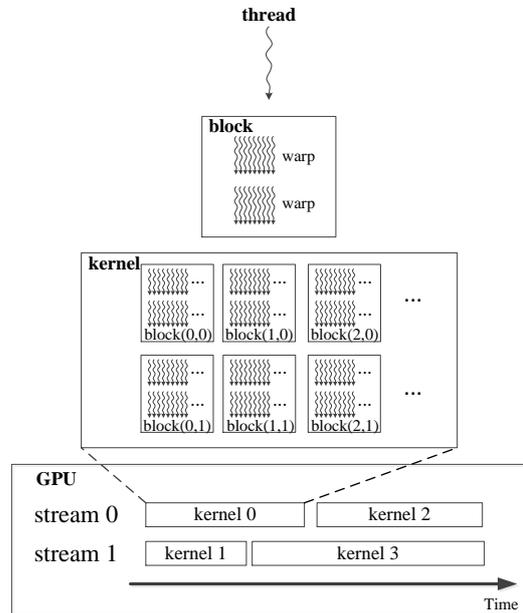


Figure 1. The hierarchy of stream, kernel, block, warp and thread.

120 method uses OpenACC directive as hints in the original CPU code as in Zhenya et al. (2010). The last method is the most complex but also the most effective; it involves rewriting the entire program using low level CUDA subroutines.

In CUDA terminology, a kernel is a single section of code or subroutine running on the GPU. The underlying code in a kernel is split into a series of threads each of which deals with different
 125 data. These threads are grouped into equal size thread blocks that can be executed independently. A thread block is further divided into warps as basic scheduled units. A warp consists of 32 consecutive threads that execute the same instruction simultaneously. Each kernel and data transfer command in CUDA has an optional parameter “stream ID”. If the “stream ID” is set in code, commands belonging to different streams can be executed concurrently. A stream in CUDA is a sequence of
 130 commands executed in order. Different streams can execute concurrently with different priorities. Figure 1 illustrates the hierarchy of these terms.

At present, CUDA compilers are available for C and Fortran. Although CUDA-Fortran has been available since 2009 and would involve less modification of the mpiPOM code, we chose CUDA-C to convert the POM.gpu-v1.0 because of the following reasons: 1) CUDA-C is free of charge; 2)
 135 previous work (Henderson et al., 2011) has shown that the CUDA-Fortran compiler did not perform as well as the CUDA-C version for some of the kernels during the porting of NIM; 3) the read-only data cache is not supported by CUDA-Fortran, which is the key optimization of Section 4.1.2; and 4) we have many previous optimisation experiences with CUDA-C.

4 Full GPU acceleration of the mpiPOM

140 Figure 2 is a flowchart illustrating the structure of POM.gpu. The main difference between the mpiPOM and the POM.gpu is that the CPU in the POM.gpu is only responsible for the initializing and the output work. The POM.gpu begins by initializing the relevant arrays on the CPU and then copies data from CPU to the GPU. The GPU then performs all of the model computations. Outputs such as velocity and sea surface height, are copied back to the CPU and are then written to
145 the disk at a user-specified time interval.

In the following sections, we introduce the optimizations of the POM.gpu by computation, communication and I/O aspects individually.

For the individual GPUs, we concentrate on memory access optimization by making better use of caches in the GPU memory hierarchy. This involves using read-only data cache, local memory
150 blocking, loop fusion, function fusion, and disabling error-correcting code memory. The test results demonstrate that a single GPU can run the model almost one hundred times faster than a single CPU core.

In terms of communication, we overlapped the sending of boundary data between the GPUs with the main computation. Data is also sent directly between the GPUs, bypassing the CPU.

155 In terms of I/O, we launched extra MPI processes on the main CPU to output the data. These MPI processes are divided into two categories, the computation processes and the I/O processes. The computation processes are responsible for launching kernels into GPUs and the I/O processes are responsible for copying data back from the GPUs and for writing to disks. The computation processes and the I/O processes can execute simultaneously to save output time.

160 4.1 Computational optimizations in a single GPU

Managing the significant performance difference between global memory and on-chip fast memory is the primary concern for GPU computing. The ratio of bandwidth between global memory and shared memory is approximately 1:10. Therefore, data reuse in on-chip cache always needs to be seriously considered. As shown on the right side of Fig. 3, we propose two classes of optimization,
165 including the standard optimization of fusion and the special optimization of the GPU, to better utilize the fast registers and caches.

4.1.1 Standard optimizations of fusion

Fusion optimization in the POM.gpu code includes loop fusion and function fusion. The loop fusion merges several loops into one loop and the function fusion merges several subroutines into one
170 subroutine.

Loop fusion is an effective method to store scalar variables in registers for data reuse. As shown in Fig. 4, if the variable $drho_x(k, j, i)$ is read several times in multiple loops, we can fuse these loops

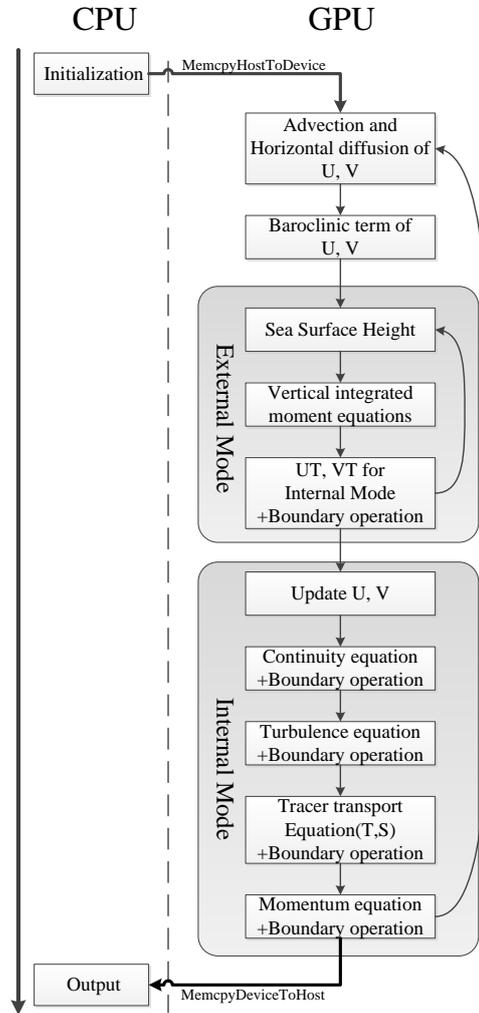


Figure 2. POM.gpu flowchart

into one. Therefore, $drho_x(k, j, i)$ will first be read from the global memory and then repeatedly read from a register. For instance, for the *profq* kernel optimized with loop fusion, the device memory transactions decrease by 57%, and the running speed of this kernel is improved by 28.6%. The loop fusion optimization can also be applied in a number of mpiPOM subroutines.

Similar to loop fusion, we can also merge functions in which the same arrays are accessed. For example, the *advv* and *advu* functions in the mpiPOM code are used to calculate the advection terms in horizontal directions, respectively. After merging them into one subroutine, the redundant memory access is avoided. The function fusion can also be applied in which one function is called several times to calculate different tracers. The *proft* function in the mpiPOM code is called twice

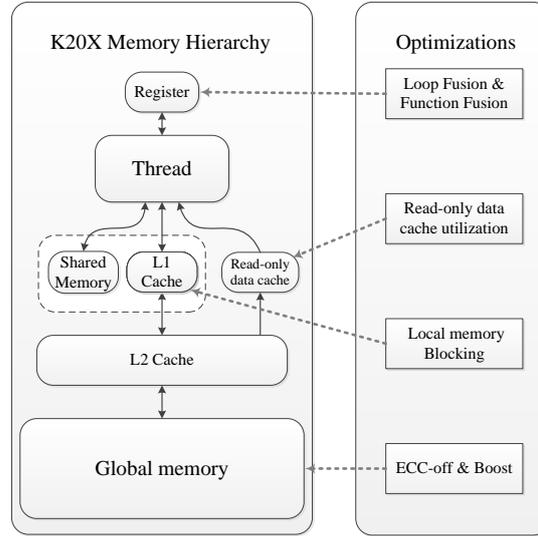


Figure 3. The memory hierarchy of the K20X GPU and the relationships with each optimization

```

/*****
*There exist two loops.
*drhox is visited twice in these loops.
*****/
for (k = 1; k < nz-1; k++){
    drhox[k][j][i] = drhox[k-1][j][i] + A[k][j][i];
}

for (k = 0; k < nz-1; k++){
    drhox[k][j][i] = drhox[k][j][i] * B[k][j][i];
}

```

(a)Original CUDA-C code

```

/*****
*These loops can be fused into one
*to reduce global memory access.
*****/
for (k = 1; k < nz-1; k++){
    drhox[k][j][i] = drhox[k-1][j][i] + A[k][j][i];
    drhox[k-1][j][i] = drhox[k-1][j][i] * B[k-1][j][i];
}

drhox[k-1][j][i] = drhox[k-1][j][i] * B[k-1][j][i];

```

(b)Optimized CUDA-C code

Figure 4. A simple example of loop fusion.

– one for temperature and one for salinity. Their computing formulas are similar and some common arrays are accessed. After function fusion, the running speed of the *profit* kernel is improved by 28.8%.

185 4.1.2 Special optimizations of the GPU

Our special optimizations mainly focus on the improved utilization of the read-only data cache and the L1 cache on the GPU. It is useful to alleviate the bottleneck of memory bandwidth that is limited by using these fast on-chip caches.

```

/*****
*3D arrays ee and gg represent row transformation
*coefficients of the chasing method.
*****/
for (k = 1; k < nz-2; k++){
    ee[k][j][i] = ee[k-1][j][i]*A[k][j][i];
    gg[k][j][i] = ee[k-1][j][i]*gg[k-1][j][i] - B[k][j][i];
}

for (k = nz-3; k >= 0; k++){
    uf[k][j][i] = (ee[k][j][i]*uf[k+1][j][i]+gg[k]) * C[k][j][i];
}

/*****
*Each thread pulls its own tile of ee,gg to
*1D new arrays ee_new, gg_new(local memory).
*There two new arrays can be cached in L1 for reuse.
*****/
for (k = 1; k < nz-2; k++){
    ee_new[k] = ee_new[k-1]*A[k][j][i];
    gg_new[k] = ee_new[k-1]*gg[k-1] - B[k][j][i];
}

for (k = nz-3; k >= 0; k++){
    uf[k][j][i] = (ee_new[k]*uf[k+1]+gg_new[k])*C[k][j][i];
}

```

(a)Original CUDA-C code

(b)Optimized CUDA-C code

Figure 5. A simple example of local memory blocking.

There is a 48KB read-only data cache in the K20X GPU. We can automatically use this as long
190 as the read-only condition is met. In the POM.gpu, we simply add “const __restrict__” qualifiers
into the parameter pointers to explicitly direct the compiler to implement the optimization. As an ex-
ample, consider the calculations of advection and the horizontal diffusion terms. Because mpiPOM
adopts the Arakawa C-grid, in the horizontal plane, updating the temperature(T) requires the veloc-
ity of longitude(u), the velocity of latitude(v) and the horizontal kinematic viscosity(aam) on the
195 neighbouring grid points. In one kernel, the u and v arrays are accessed twice, and the aam array is
accessed four times. After using read-only data cache to improve the data locality, the running speed
of this kernel is improved by 18.8%.

To reuse the data in each thread, we use local memory blocking to pull the data from global
memory to the L1 cache. In this method, a small subset of a dataset is loaded into the fast on-
200 chip memory and then the small data block is repeatedly accessed by the program. This method is
helpful in reducing the need to access the off-chip with high latency memory. In the subroutines of
the vertical diffusion and source/sink terms, the chasing method is used to solve a tridiagonal matrix
along the vertical direction for each grid point individually. Each thread only accesses its own tiles of
row transformation coefficients. As shown in Fig. 5, the arrays are accessed twice within one thread,
205 one from the surface($k = 0$) to the bottom($k = nz - 1$) and another from the bottom($k = nz - 1$) to
the surface($k = 0$). After blocking the vertical direction arrays in local memory, the L1 cache is fully
utilized and the running speed of these subroutines is improved by 35.3%.

In current implementation, as in the original mpiPOM code, the 3D arrays of variables are stored
sequentially as east-west(x), north-south(y), vertical(z), i.e., i, j, k ordering. 2D arrays are stored in
210 i, j ordering. The vertical diffusion is solved using a tridiagonal solver which is calculated sequen-
tially in the z direction. For simplicity, in our kernel functions the grid is divided along x and y .
Each GPU thread then specifies an (x, y) point in the horizontal direction and performs all of the

calculations from the surface to the bottom. The thread blocks are divided as 32×4 subdomains in the x - y plane. In the x direction, the block number must be a multiple of 32 threads to perform consecutive and aligned memory access within a warp (NVIDIA, 2015). In the y direction, we tested many thread numbers, such as 4 and 8, and obtained similar performances. We ultimately choose 4 because this value produced more blocks and allowed us to distribute the workload more uniformly amongst the SMs. In addition, $128(=32 \times 4)$ threads are enough to maintain the full occupancy, which is the number of active threads in each multiprocessor.

In GPU computing, one is free to choose which arrays will be stored in an on-chip cache. Our experience involves putting the data along the horizontal direction into the read-only cache to reuse among threads, and putting the data along with vertical direction into the local memory for reuse within one thread.

Furthermore, we improve the global memory bandwidth by disabling the Error Checking and memory Correcting(ECC-off), as well as enhancing the clock on the GPU(GPU boost). This method improves the performance of the POM.gpu by 13.8%.

4.1.3 Results of the computational optimizations

We divide all of the POM.gpu subroutines into three categories based on their different computational patterns. As shown in Table 2, in the POM.gpu, we deploy different optimizations in these categories to improve the performance of POM.gpu; these categories are described as follows.

(1) Category 1: Advection and horizontal diffusion (*adv*)

This category has 6 subroutines, and calculates the advection, horizontal diffusion and the pressure gradient and Coriolis terms in the case of velocity. Here, it is possible to reuse data among adjacent threads, and the subroutines therefore benefit from using the read-only data cache. At the same time, the variables are calculated in different loops or in different functions such that the loop fusion and function fusion optimizations are applied to this part as well.

(2) Category 2: vertical diffusion (*ver*)

This category has 4 subroutines, and calculates the vertical diffusion. In this part, the chasing method is used in the tridiagonal solver in the k -direction. The main feature is that the data are accessed twice within one thread, once from the surface to the bottom and again from the bottom to the surface. The subroutines are significantly sped up after grouping the k -direction variable in the local memories.

(3) Category 3: vorticity (*vort*), baroclinicity (*baro*), continuity equation (*cont*) and equation of state (*state*)

This category is less time consuming than the two categories described above, but it also benefits from our optimizations. Because data exists reuse among threads, the use of a read-only data cache improves data locality. For the *vort* subroutine, there is data reuse within one thread, and thus the loop fusion improves the data locality.

Table 2. Different subroutines adopt different optimizations in the POM.gpu

Subroutines	Loop fusion	Function fusion	Read-only data cache	Local memory blocking	ECC-off & GPU boost	Speedup
Adv & Hor diff	✓	✓	✓		✓	2.05X
Ver diff	✓	✓		✓	✓	2.82X
Baroclinicity	✓		✓		✓	2.08X
Continuity equation			✓		✓	1.39X
Vorticity	✓		✓		✓	3.19X
State equation			✓		✓	1.35X

4.2 Communication optimizations among multiple GPUs

250 In this section, we present the optimizing strategies for multiple GPUs. In the mpiPOM, the entire domain is split along the horizontal directions and each MPI process is responsible for the model’s computation of one subdomain, following Jordi and Wang (2012). In the POM.gpu, we attach one MPI process to one GPU and move the complete computation to the GPU. The MPI process is in charge of the computation within each subdomain and of the data transfer between the GPU and main memory. The data transfer between subdomains is handled by the GPUs directly. Shimokawabe et al. (2010) and Yang et al. (2013) proposed fine-grained overlapping methods of GPU computation and CPU communication to improve the computing performance. An important issue in their work is that the communications between multiple GPUs explicitly require the participation of the CPU. In our current work, we simply bypass the CPU in implementing the communication to fully exploit the capability of the GPUs.

At present, two MPI libraries, OpenMPI and MVAPICH2, provided support for the direct communication from the GPU to the main memory. This capability is referred as CUDA-aware MPI. We attempted to use MVAPICH2 to implement direct communication among multiple GPUs. However, we found that inter-domain communication occupied nearly 18% of the total runtime.

265 Instead, to fully overlap the boundary operations and MPI communications with computation, we adopt the data decomposition method shown in Fig. 6. The data region is decomposed into three regions: the inner region, the outer region, and a halo region which exchanges data with its neighbours. In our design, the inner region, which is the most time-consuming is allocated to “stream 1”. The east/west outer region is allocated to “stream 2” and the north/south outer region is allocated to “stream 3”. In the east/west outer region, the width is set to 32 to ensure consecutive and aligned memory access in a warp. All of the halo regions are also allocated to “stream 2”.

The workflow of multiple streams on the GPU is shown in Fig. 7. The east/west/north/south regions are common kernel functions that can run in parallel with the inner region through different streams. The communication operations between domains are implemented by an asynchronous

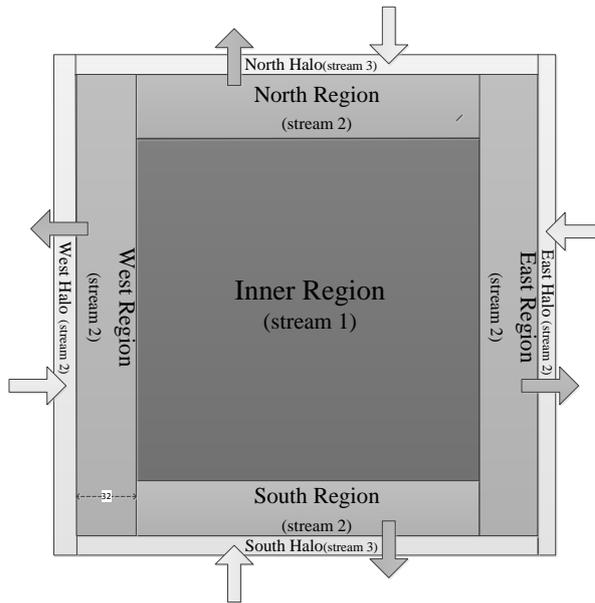


Figure 6. Data decomposition in the POM.gpu

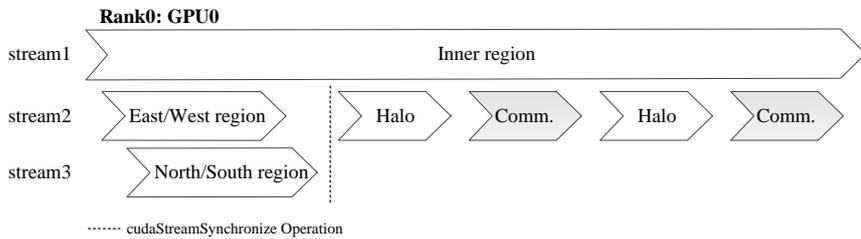


Figure 7. The workflow of multiple streams on the GPU. The “inner/east/west/north/south region” and “Halo” refer to the computation and update of the corresponding region. “Comm.” refers to the communication between processes, which implies synchronization.

275 CUDA memory copy. The corresponding synchronization operations between the CPU and the GPU or among the MPI processes are implemented by a synchronization CUDA function and a MPI barrier function. To overlap the subsequent communication with the inner region, “stream 2” and “stream 3” for the outer region have higher priority in preempting the computing resource from “stream 1” at any time. Based on this workflow, the inter-domain communication is overlapped with
 280 the computation. The experimental results show that our design can remove the communication overhead taken by MVAPICH2.

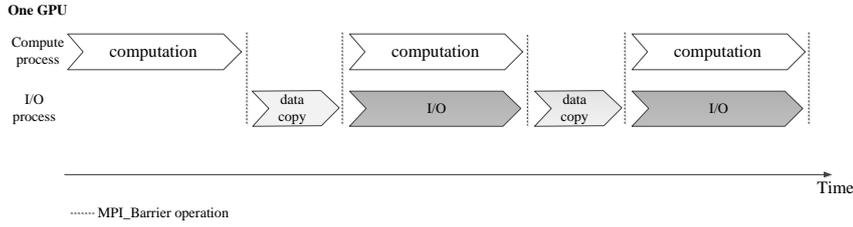


Figure 8. One computing process and one I/O process both set their contexts on the same GPU. During the data copy phase, the computing process remains idle and the I/O process will copy data from the GPU to the CPU through the *cudaMemcpy* function.

4.3 I/O optimizations between the GPUs and the CPUs

The time consumed for I/O in the mpiPOM is not significant. However, after we fully accelerate the model by GPU, it accounts for approximately 30% of the total runtime. The computing phase and the I/O phase are serial, which means that the GPU will remain idle until the CPU finishes the I/O operations. Motivated by previous work on I/O overlapping (Huang et al., 2014), we designed a similar method following computations on a GPU and I/O operations on a CPU to run in parallel.

In the POM.gpu, we chose to launch more MPI processes. The MPI processes are divided into computing processes and I/O processes with different MPI communicators. The computing processes are responsible for launching kernel functions as usual, and the I/O processes are responsible for output. One I/O process attaches to one computing process and these two processes set their contexts on the same GPU.

Because the I/O processes must fetch data from the GPU, communication is necessary between them. The I/O processes obtain the device buffer pointers from the computing processes during the initialization phase. When writing history files, the computing processes are blocked and remain idle for a short time, waiting for I/O processes to fetch data. Then, the computing processes continue their computation, and the I/O processes complete their output in the background, as illustrated in Fig. 8. This method can be further optimized by placing the archive data into a set-aside buffer and carry on the main calculation. However, the method requires more memory, which is not abundant in current K20X GPUs.

The advantage of this method is that it overlaps the I/O on the CPU with the model calculation on the GPU. In serial I/O, the GPU computing processes are blocked while data are sent to the CPU and written to disk. In overlapping I/O, the computing processes only wait for the data to be sent to the host. The bandwidth of data brought to the host is approximately 6 GB/s, but the output bandwidth to the disk is approximately 100 MB/s, as determined by the speed of the disk. Therefore, the overlapping method significantly accelerates the entire application.

5 Experiments

In this section, we first describe the specification of our platform and comparison methodology to validate the correctness of the POM.gpu. Furthermore, we present the performance and scalability
310 of the POM.gpu compared with the mpiPOM.

5.1 Platform Setup

The POM.gpu runs in a workstation consisting of two CPUs and four GPUs. The CPUs are 2.6 GHz 8-core Intel SandyBridge E5-2670. The GPUs are Nvidia Tesla K20X. The operating system is RedHat Enterprise Linux 6.3 x86_64. All programs are compiled with Intel compiler v14.0.1,
315 CUDA 5.5 Toolkit, Intel MPI Library v4.1.3 and MVAPICH2 v1.9.

For comparison, the mpiPOM runs on the *Tansuo*100 cluster at Tsinghua University consisting of 740 nodes. Each node is equipped with two 2.93 GHz 6-core Intel Xeon X5670 CPUs and 32 GB of memory. The nodes are connected through an infiniband network. The operating system is RedHat Enterprise Linux 5.5 x86_64. Programs on this platform are compiled with Intel compiler
320 v11.1 and Intel MPI v4.0.2. The mpiPOM code is compiled with its original compiler flags, i.e., “-O3 -fp-model precise”.

5.2 The test case and the verification of accuracy

The “dam break” simulation (Oey, 2014) is conducted to verify the correctness and test the performance and scalability of the POM.gpu. It is a baroclinic instability problem that simulates flows
325 produced by horizontal temperature gradients. The model domain is configured as a straight channel with a uniform depth of 50 m. Periodic boundary conditions are used in the east-west direction, and the channel is closed in the north and south. Its horizontal resolution is $1\text{km} \times 1\text{km}$. The domain size of this test case is 962×722 horizontal grid points and 51 vertical sigma levels, which is limited by the capacity of one GPU’s memory. Initially, the temperature in the southern half of the channel
330 is 15°C and 25°C in the northern half. The salinity is fixed at 35 psu. The fluid is then allowed to adjust. In the first 3-5 days, geostrophic adjustments occur. Then, an unstable wave develops due to baroclinic instability. Eventually, eddies are generated. Figure 9 shows the sea-surface height (SSH), sea-surface temperature (SST), and currents after 39 days. The scales of the frontal wave and eddies are determined by the Rossby radius of deformation. This dam break case uses a single-precision
335 format.

To verify the accuracy, we check the binary output files of the mpiPOM and the POM.gpu, as in Mak et al. (2011). The test results demonstrate that the variables velocity, temperature, salinity and sea surface height are all identical.

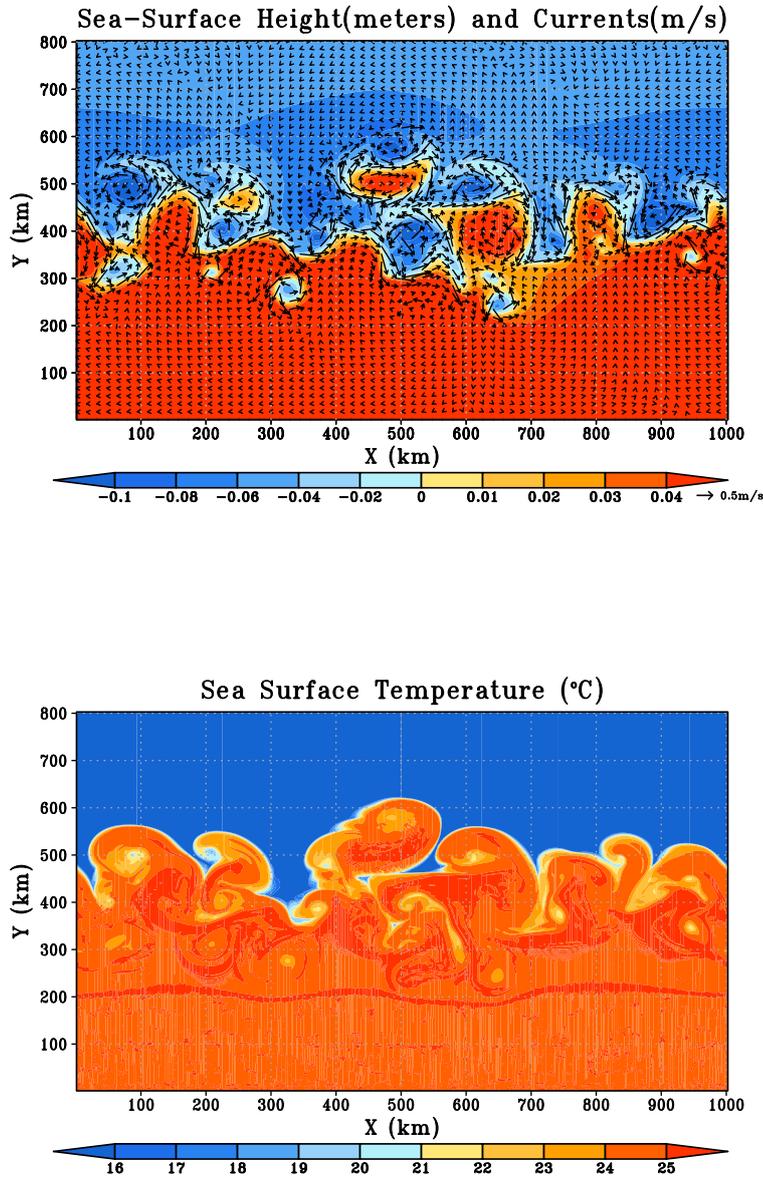


Figure 9. The model results after 39 days of simulation. For the top figure, the colour shading is the sea-surface height (SSH), and vectors are ocean currents. For the bottom figure, the colour shading is the sea surface temperature (SST). Several warm and cold eddies are generated in the middle of the domain where the SST gradient is largest; their scales are determined by the Rossby radius of deformation.

5.3 Model Performance

340 To understand the advantages of the optimizations in Sec. 4, we conducted different tests. The metrics of seconds per simulation day are measured to compare the model performance.

5.3.1 Single GPU performance

In our first test, we compare the performance of the mpiPOM using two different CPUs, the Intel X5670 CPU (6 cores) and the Intel E5-2670 CPU (8 cores), with that obtained from the POM.gpu using one single GPU. Fig. 10 shows that one K20X GPU can compete with approximately 55 E5-2670 CPU cores to 95 X5670 CPU cores in the simulation. From the parameters of the Intel E5-2670 CPU and Nvidia K20X GPU, we find that, the ratio of memory bandwidth and the ratio of floating points performance are approximately 1:5 and 1:10, respectively. This means, if an application is strictly memory bandwidth limited, one GPU can compete with 5 CPUs; if an application is strictly computation limited, it can compete with 10 CPUs. Since the mpiPOM is memory bandwidth limited, the POM.gpu should provide equivalent performance to the mpiPOM running on up to $5 \times 8 = 40$ CPU cores. Our procedure attempts to optimize memory access and we can further increase this number to 55.

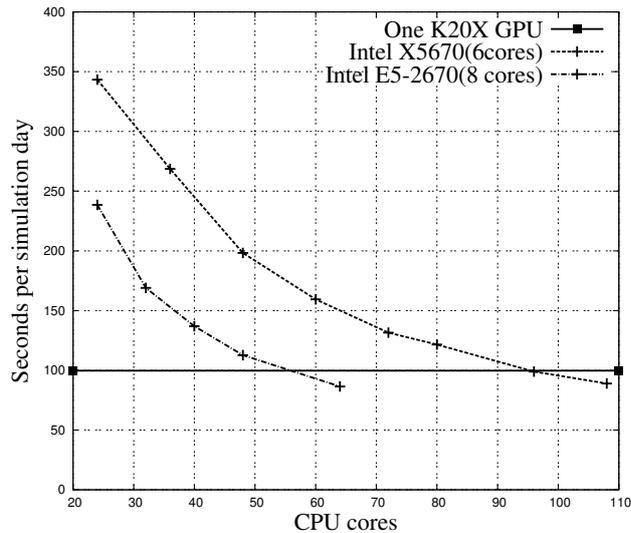


Figure 10. Performance comparison with different hardware platforms

5.3.2 Multiple GPUs performance

In the second test, we compare our communication overlapping method with the MVAPICH2 library. Fig. 11 presents the weak scaling performance on multiple GPUs, where the grid size for each GPU is kept at $962 \times 722 \times 51$. When 4 GPUs are used with MVAPICH2, approximately 18% of the total runtime is consumed by inter-domain communication and boundary operations. This overhead can be greatly reduced by our communication overlapping method.

In the third test, we fix the global grid size at $962 \times 722 \times 51$, and measure the strong scaling performance of POM.gpu. Table 3 shows that the strong scaling efficiency is 99% on 2 GPUs and 92% on 4 GPUs. When more GPUs are used, the size of each subdomain becomes smaller. This

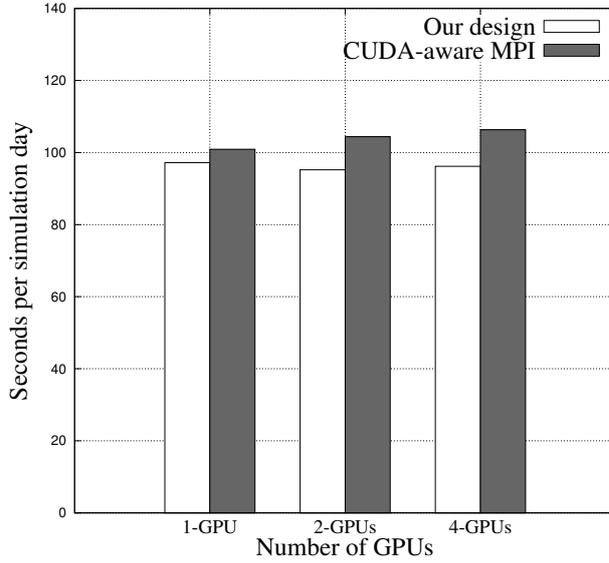


Figure 11. The weak scaling test between our communication overlapping method and the MVAPICH2 sub-routines.

decreases the performance of POM.gpu in two aspects. First, the communication overhead may exceed the computation time of the inner region as the size of each subdomain decreases. As a result, the overlapping method in Section 4.2 are not effective. Second, there are many “small” kernels in the POM.gpu code, in which the calculation is simple and less time consuming. With fewer inner region computations, the overhead of kernel launching and implicit synchronization with kernel execution must be counted.

Table 3. The strong scaling result of POM.gpu

Number of GPUs	1-GPU	2-GPUs	4-GPUs
Time(s)	97.2	48.7	26.3
Efficiency	100%	99%	92%

5.3.3 I/O performance

In the fourth test, we compare our I/O overlapping method with the parallel NetCDF (PnetCDF) method and NO-I/O. NO-I/O means that all I/O operations are disabled in the program and that the time measured is the pure computing time. This simulation is run for 20 days, and the history files are output daily. The final history files in NetCDF format are approximately 12 GB. Fig. 12 shows that the I/O overlapping method outperforms the PnetCDF method. For 1 GPU and 2 GPUs, the overall runtime decreases from 1694/1142 seconds to 1239/688 seconds, which is close to the NO-

I/O. The extra overhead of our method compared with NO-I/O involves the computing processes that need to be blocked until the I/O processes obtain data from the GPUs. When running with 4 GPUs, the output time exceeds the computation time. Then, the I/O phase cannot be fully overlapped with the model computation phase. The overall runtime equals the sum of the computation time and the non-overlapped I/O time.

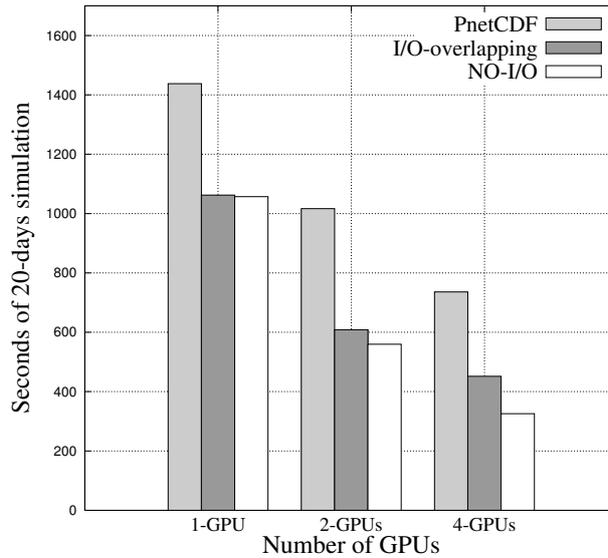


Figure 12. I/O test for the POM.gpu

5.3.4 Comparison with a cluster

In the last test, we compare the performance of POM.gpu on a workstation containing 4 GPUs with that on the *Tansuo100* cluster. Three different high-resolution grids (Grid-1: $962 \times 722 \times 51$, Grid-2: $1922 \times 722 \times 51$, Grid-3: $1922 \times 1442 \times 51$) are used. Fig. 13 shows that our workstation with 4 GPUs is comparable to 408 standard CPU cores (= 34 nodes \times 12 cores/node) in the simulation. Because the thermal design power of one X5670 CPU is 95 W and that of one K20X GPU is 235 W, we reduce the energy consumption by a factor of 6.8. Theoretically, as the subdomain of each MPI process becomes smaller, the cache hit ratio of the mpiPOM code will increase. This will greatly alleviate the memory bandwidth-limited problem. However, in the simulation on 408 standard CPU cores, the MPI communication may occupy more than 40% of the total execution time. When scaling to over 450 cores, the mpiPOM simulation may instead become slower, as shown in Fig. 13. Therefore, for high-resolution ocean modelling, our POM.gpu has a clear advantage compared to the original mpiPOM.

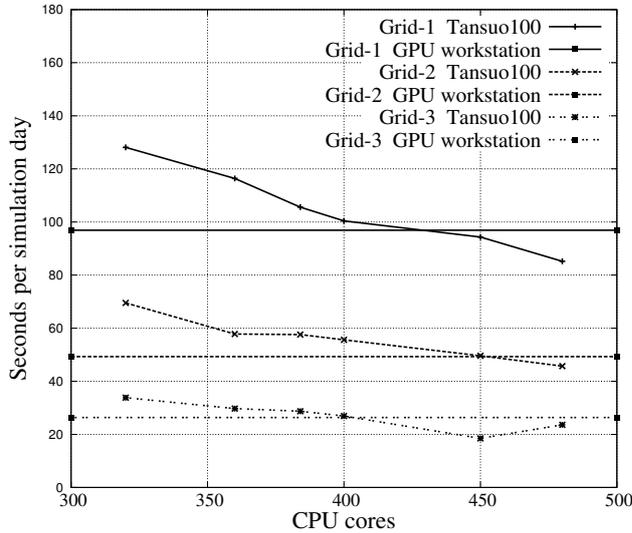


Figure 13. Performance test of four GPUs compared with *Tansuo100* cluster

6 Code availability

395 The POM.gpu version 1.0 is available at <https://github.com/hxmhuang/POM.gpu>. To reproduce the test case in Section 5, the script “run_exp002.sh” is provided to compile and execute the POM.gpu code.

7 Conclusions and future work

In this paper, we develop POM.gpu, a full GPU solution based on the mpiPOM. Unlike previous
 400 GPU porting, the POM.gpu code distributes the model computations on the GPU. Our main contributions include: optimizing the code on each of the GPUs, the communications between GPUs, and the I/O process between the GPUs and the CPUs. Using a workstation with 4 GPUs, we achieve the performance of a powerful CPU cluster with 408 standard CPU cores. Our model also reduces the energy consumption by a factor of 6.8. It is a cost-effective and energy-efficient strategy for
 405 high-resolution ocean modelling. We have described the method and tests in details and, with the availability of the POM.gpu code, our experiences may hopefully be useful to developers and designers of other general circulation models.

In our current POM.gpu, we design a large number of kernel functions because we port the entire mpiPOM one subroutine at a time. This was done to simplify the debugging of POM.gpu and to
 410 check that the results are consistent with the mpiPOM. In our future work, we will adjust the code structure of POM.gpu and adopt aggressive function fusion to further improve the performance.

Previous studies proposed to take advantage of data locality between time steps by time skewing (McCalpin and Wonnacott, 1999; Wonnacott, 2000), thus transforming the problem of memory

bandwidth into the problem of computation. However, the real-world ocean models, including the
415 mpiPOM, often involve hundreds of thousands lines of code, and analysing the data dependency
and applying time skewing in such a context are challenging and difficult. We leave that to the next-
generation POM.gpu.

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